

To infinity – and beyond!

Julian Bradfield

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# Infinity and eternity

for ever and ever (idiom)

immer und ewig (idiom)

for ever and a day (Shakespeare)

nunc et semper et in saecula saeculorum (from Gk, probably from Aramaic idiom)

*The King said, "The third question is, how many seconds of time are there in eternity." Then said the shepherd boy, "In Lower Pomerania is the Diamond Mountain, which is a league high, a league wide, and a league in depth; every hundred years a little bird comes and sharpens its beak on it, and when the whole mountain is worn away by this, then the first second of eternity will be over." (from Grimm)*

## Towards infinity

*Mathematicians and other children often play the following game: We take turns naming numbers, and see who can name the largest one. This is a game in the psychological rather than the formal sense, since I might always just add one to your number, but my goal is to try to completely demolish your ego by transcending your number via some completely new principle.*

Kenneth Kunen

# Georg Cantor

'The father of set theory'

1845–1918

Martin-Luther-Universität  
Halle-Wittenberg

1874: the birth of set  
theory, and the discovery  
of different levels of  
infinity

1883: the theory of ordinal  
numbers



# Ordinals and Cardinals

Language:

- ▶ **cardinal** numerals
  - ▶ *one, two, ...*
  - ▶ “how many?”
- ▶ **ordinal** numerals
  - ▶ *first, second, ...*
  - ▶ “where in a sequence?”

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## Language:

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  - ▶ “how many?”
- ▶ **ordinal** numerals
  - ▶ *first, second, ...*
  - ▶ “where in a sequence?”

## Mathematics:

- ▶ **cardinal** numbers
  - ▶ *0, 1, 2, ...*
  - ▶ “how many [in a set]?”
- ▶ **ordinal** numbers
  - ▶ *0, 1, 2, ...*
  - ▶ “where in a sequence?”, also “how long [is a sequence]?”

The sequence  $a,b,a$  is 3 letters long, but contains 2 distinct letters.

# Counting

0:

1: |

2: ||

10: |||||

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Obviously we can keep counting 'for ever':

|||||



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$\omega$ : ||||| ...

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$\omega + 1$ : ||||| ... |

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$\omega$ : ||||| ...

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$\omega + 1$ : ||||| ... |

Write  $\omega$  for the length of the infinite sequence.

To help visualization, compress the infinite sequence to

|||.

# Addition of ordinals

Adding sequences is just putting one after the other:

$\omega + 1$ : |||||

$\omega + 3$ : |||||

$\omega + \omega$ : |||||

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$\omega + 1$ :  $\|\|_n\|$

$\omega + 3$ :  $\|\|_n\|\|\|$

$\omega + \omega$ :  $\|\|_n\|\|_n$

But  $1 + \omega$ :  $\|\|_n = \omega$

# Multiplication of ordinals

Integer multiplication is just repeated addition:  $2 \times 3 = 2 + 2 + 2$ .

By convention, let's write  $x \cdot y$  to mean  $y$  copies of  $x$  added together.

$2 \cdot 3$ : || || ||

$\omega \cdot 3$ : |||, |||, |||

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which we might visualize as



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Key property of these ordinals: they are *well-founded*. Start with an infinite ordinal, and keep decreasing it: after a **finite** (but arbitrarily large) number of steps, you must hit zero.



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This means that ordinals can generalize *proof by induction*:

If

- ▶  $P(\alpha)$  holds for  $\alpha = 0$ , and
- ▶ **if**  $P(\beta)$  holds for all  $\beta < \alpha$ , **then**  $P(\alpha)$  holds,

then we can conclude that  $P(\alpha)$  holds for all ordinals  $\alpha$ .

(Often restricted to ordinals less than some fixed  $\alpha_0$ .)

## Example: why does Ackermann terminate?

The *Ackermann* function (of two integer arguments)  $A(x, y)$  is defined recursively thus:

$$A(0, y) = y + 1$$

$$A(x, 0) = A(x - 1, 1) \quad \text{for } x > 0$$

$$A(x, y) = A(x - 1, A(x, y - 1)) \quad \text{for } x, y > 0$$

Is it obvious that this recursive computation ever finishes on, e.g.,  $A(4, 4)$ ?

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In each recursive call, *either*  $x$  gets smaller, *or*  $x$  stays the same and  $y$  gets smaller.

This is an induction on  $\omega \cdot \omega$ .

The Ackermann function grows quite fast – see later ...

# Ordinal Exponentiation

Integer exponentiation is just repeated multiplication:

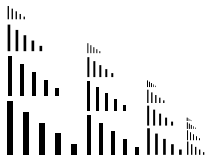
$$2^3 = 2 \times 2 \times 2.$$

I.e., we write  $x^y$  for  $y$  copies of  $x$  multiplied together.

$$\omega^2: \omega \cdot \omega$$

$$2^\omega: 2 \cdot 2 \cdot 2 \cdot \dots = \omega$$

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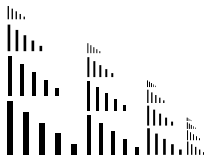
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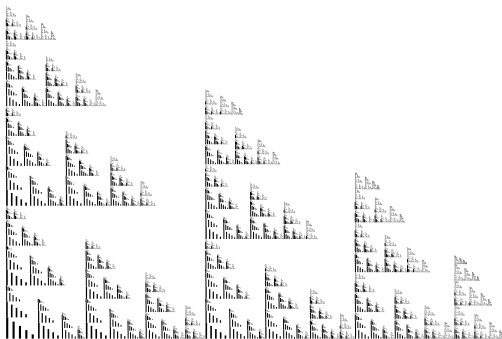
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$$\omega^\omega: \omega \cdot \omega \cdot \omega \cdot \dots$$



## A little puzzle . . .

A number is written in *hereditary base*  $b$  if it's a sum of powers of  $b$ , with all the exponents also written in hereditary base  $b$ . E.g. with  $b = 2$

$$1030 = 2^{10} + 2^2 + 2 = 2^{2^{2+1}+2} + 2^2 + 2$$

or with  $b = 3$

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Think of a number  $n$ . Write it in h.b. 2. Now replace 2 by 3 and evaluate. Subtract 1. Write the result in h.b. 3. Replace 3 by 4 and evaluate. Subtract 1. And so on ... until you hit zero. Let  $G(n)$  be the length of this process – if it finishes!

For example:  $G(3)$

$$\begin{aligned} 3 &= 2 + 1 \\ &\rightarrow 3 + 1 = 4 \\ 4 - 1 &= 3 = 3 \\ &\rightarrow 4 = 4 \\ 4 - 1 &= 3 = 4 \\ &\rightarrow 3 = 5 \\ 3 - 1 &= 2 = 5 \\ &\rightarrow 2 = 6 \\ 2 - 1 &= 1 = 6 \\ &\rightarrow 1 = 7 \\ 1 - 1 &= 0 \end{aligned}$$

So  $G(3) = 6$



$G(4) =$

68950808030926201657363899596115099569577498758029736589  
65164942362743495979724871888253075446277672715412687741  
34196294274754024623945165423420847416977379911463833552  
69129320073235045130731133415321473276443100557449932505  
15006661770697335697266822986380629230539311939473732984  
32189645058087369473177341975229512418408401173732994662  
36583517126642762404390343968364036246706786021125426974  
22457548590135058038973996050222167215602290558339854433  
64582849621578912386681708820717886170299010486094304298  
31938313300623537993032219144244347215613123094143176938  
67571586750377935644459245645595556087522305546773436198  
47032332425407785083961078958596196387897297104581844575  
77157677751206967346327625413465613506947655384380307508  
65233130216216163628621847422406626611936799943353915562  
43559138950800725078317787807770746695975268954544726471  
50035241859874391011058882911482099143475541850986185545

We now skip 135278 slides ...

30187441025940056292813034386280177679544034517464936619  
31073504178981468557969472279965424657699404235005937914  
72930883993663027161466688743717253581936410274739801296  
83060714286243899305063868650864112105238061406944808189  
08304913462509086531545100380965533413343423478836091833  
53220182680722735478679352859535040769913825815484931187  
10726329608316620305483302616305150324000876272357296528  
10182401729583610978448023254665651115973448118179302336  
79234929512268465106495927833854484067484182464486747555  
62975216019453924341023727286959093404563639409013246678  
20328593203290715635768149137536972887886088038810819894  
08291784060416318863529224353808259669206267357619658951  
446422310193135419323844928197722374143

Or, more comprehensibly, about  $7 \times 10^{121210694}$ ; or about  $2^{2^{29}}$ .

$$G(5) =$$

$$10^{10^{\dots 10}}$$

$$G(5) =$$

$$10^{10^{\dots^{10}}} \left. \vphantom{10^{10^{\dots^{10}}}} \right\} 10^{10^{\dots^{10}}}$$

$G(5) =$

$$\left. \begin{matrix} 10 \\ 10^{10} \\ \dots \\ 10 \end{matrix} \right\} \left. \begin{matrix} 10 \\ 10^{10} \\ \dots \\ 10 \end{matrix} \right\} \begin{matrix} 10 \\ 10^{10} \\ \dots \\ 10 \end{matrix}$$

$$G(5) =$$

$$\left. 10^{10^{\dots 10}} \right\} \left. 10^{10^{\dots 10}} \right\} \left. 10^{10^{\dots 10}} \right\} 10^{10^{10^{21}}}$$

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or to put it in binary,

$$\left. 2^{2^{\dots 2}} \right\} \left. 2^{2^{\dots 2}} \right\} \left. 2^{2^{\dots 2}} \right\} 2^{2^{2^{2^6}}}$$



## Why does $G$ always terminate?

A slight variation of the description:

Think of a number  $n$ . Write it in h.b.  $2$ , and replace  $2$  by  $\omega$ ; let  $b = 2$ . Increment  $b$ , and subtract  $1$ , expanding  $\omega$  to  $b$  (only) when necessary; repeat until zero.

$$4 = \omega^\omega \qquad b = 2$$

$$\rightarrow 26 = \omega^\omega - 1 \qquad b = 3$$

$$= \omega^3 - 1 \qquad b = 3$$

$$= \omega^2 \cdot 2 + \omega \cdot 2 + 2 \qquad b = 3$$

$$\rightarrow 41 = \omega^2 \cdot 2 + \omega \cdot 2 + 1 \qquad b = 4$$

$$\rightarrow 60 = \omega^2 \cdot 2 + \omega \cdot 2 \qquad b = 5$$

$$\rightarrow 83 = \omega^2 \cdot 2 + \omega + 5 \qquad b = 6$$

The ordinal always decreases, even while its evaluation with  $\omega = b$  is increasing. This is ordinal induction.

## Fast-growing functions ...

$G(n)$  grows fast. So also does Ackermann  $A(x, y)$ :

$x \backslash y$	0	1	2	3	4
0	1	2	3	4	5
1	2	3	4	5	6
2	3	5	7	9	11
3	5	13	29	61	125

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Let  $A(n)$  mean  $A(n, n)$ . It looks as if  $A(n) > G(n)$ :

$n$	$A(n)$	$G(n)$
0	1	1
1	3	2
2	7	5
3	61	6
4	$2^{2^{65533}}$	$2^{29}$

... via iterating exponentiation ...

Multiplication  $2 \cdot n$  is iterated addition  $2 + 2 + 2 + \dots + 2$ .

Exponentiation  $2^n$  or  $2^{\wedge}n$  is iterated multiplication

$2 \times 2 \times 2 \times \dots \times 2$ .

*Tetration*  ${}^n2$  or  $2^{\wedge\wedge}n$  is iterated exponentiation  $2^{\wedge\wedge\wedge} \dots^{\wedge}2$ .

and so on ...

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Call a number *small* if it's ... small ...

... and *1-big* if it's  $2^{(\text{small})}$  ...

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Let's continue the Ackermann – Goodstein comparison:

$n$	$A(n)$	$G(n)$
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7	5-big	7-big



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5	3-big	3-big
6	4-big	5-big
7	5-big	7-big
8	6-big	$G(4)$ -big

... and so *ad infinitum*

Call a number *1-huge* if it's (1-big)-big, *2-huge* if it's (1-huge)-big, etc.

Call a number *1-humungous* if it's (1-huge)-huge, etc.  
etc. etc. etc. etc. till your brain explodes.

## ... and so *ultra infinitum*

Play the iterated exponentiation game with ordinals, and don't stop at infinity!

$$\omega, \omega^\omega, \omega^{\omega^\omega}, \dots$$

After this we need a new symbol  $\epsilon_0 = \omega^{\omega^{\omega^{\dots}}}$  – note that  $\omega^{\epsilon_0} = \epsilon_0$ , just as  $\omega^\omega = \omega \cdot \omega^\omega$  and  $\omega \cdot \omega = \omega + \omega \cdot \omega$ .

$\epsilon_0$  is the first *fixed point* of the function  $\alpha \mapsto \omega^\alpha$ .

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Then  $\epsilon_1$  is the second fixed point (and is the end of  $\epsilon_0 + 1, \omega^{\epsilon_0+1}, \omega^{\omega^{\epsilon_0+1}}, \dots$ ).

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Then there's  $\epsilon_{\epsilon_{\dots}}$ , the first fixed point of  $\alpha \mapsto \epsilon_\alpha$ .

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Then start counting fixed points of the functions that list the fixed points: the *Veblen* hierarchy. The end of this process is called  $\Gamma_0$ .

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Then it starts getting complicated ...

## Why do (some) computer scientists care?

The working theoretical computer scientist needs ordinals to do inductions. However, generally speaking, up to  $\omega^\omega$  is as far as we need to go.

*Proof theorists* (logicians, but sometimes found in CS depts!) need much bigger ordinals.

The strength of a theory (how much that is true, can it prove?) can be measured by how long are the inductions it can do.

E.g. *Primitive Recursive Arithmetic* can't do  $\omega^\omega$  inductions.

*Peano Arithmetic* can't do  $\epsilon_0$  induction – so can't prove that  $G$  terminates!

*Proof Theory* may be of interest to *Theorem Provers* ...



## Envoi

All these ordinals are small!

The real Cantorian revolution was about *cardinals* – we have not gone beyond the first infinite cardinal.

But that's for another talk.