The PEPA Workbench: User's Manual Stephen Gilmore, LFCS, Edinburgh stg@dcs.ed.ac.uk September 18, 2001

This document describes version 0.72 of the PEPA Workbench, a tool for assisting with the analysis of systems which are described by a model expressed in the stochastic process algebra PEPA. The definitive reference for PEPA is the CUP book "A Compositional Approach to Performance Modelling" by Jane Hillston (published 1996). PEPA software and accompanying documentation such as this can be obtained via the World-Wide Web from the address http://www.dcs.ed.ac.uk/pepa.

The PEPA Workbench is intended for use with a linear algebra solution tool such as those provided within the Maple and Matlab computing environments. The objective of the modelling study is to calculate the steady-state probability distribution for the system and to derive performance measures from this.

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We describe the components of the PEPA input language for the Workbench via a simple example. Consider a PEPA model with two copies of a component, P_1 , executing in a pure parallel synchronisation. P_1 is a simple sequential process which undergoes a *start* activity with rate r_1 to become P_2 which runs with rate r_2 to become P_3 which goes back to P_1 via a *stop* activity with rate r_3 .

 $\begin{array}{rcl} P_1 & \stackrel{def}{=} & (start, r_1).P_2 \\ P_2 & \stackrel{def}{=} & (run, r_2).P_3 \\ P_3 & \stackrel{def}{=} & (stop, r_3).P_1 \end{array}$

The system which we study is $P_1 \bowtie_{\emptyset} P_1$ (equivalent notation for this is $P_1 \parallel P_1$).

We have made use of aspects of the mathematical syntax for PEPA in this definition. Before solving this model we first need to encode these definitions and the pure parallel synchronisation term in the ASCII syntax accepted by the Workbench. We place these definitions in the file tiny.pepa.

```
#P1 = (start, r1).P2;
#P2 = (run, r2).P3;
#P3 = (stop, r3).P1;
P1 <> P1
```

This tiny example shows the form of all PEPA definitions which are accepted by the PEPA Workbench; a [non-empty] sequence of definitions of sequential components followed by a parallel composition of these to form the description of the system to be studied.

The purpose of the PEPA Workbench is to process models such as these in order to compute the state-space and the transitions of the model. Assuming that you access the PEPA Workbench by typing **pwb** then the model is processed as shown in Figure 1 where the user only typed the **pwb** command and the filename tiny and the rest was produced by the Workbench.

```
[unix]: pwb
PEPA Workbench Version 0.72 [maple, 22-12-1998]
Filename: tiny
Compiling the program
Writing the state table file to tiny.table
Writing the transition file to tiny.maple
The model has 9 states
The model has 18 transitions
Writing the hash table file to tiny.hash
Exiting PEPA workbench.
```

Figure 1: A sample PEPA Workbench session

You now have three files which you did not have before: tiny.table, tiny.hash and tiny.maple¹. All of these are text files which can be read on the screen or searched using a UNIX utility such as grep or printed on a printer. The real content of the analysis is in tiny.maple so we shall consider it first. It contains the entries for the transition matrix \mathbf{Q} . Depicted mathematically, \mathbf{Q} would be the matrix shown below where blank entries indicate zeroes. It does not matter for a tiny example such as this one but as a general point the transition matricies of any Markov processes are best stored within a computer as *sparse* matricies.

$$\begin{pmatrix} -2r_1 & r_1 & r_1 & & & \\ & -r_1 - r_2 & r_2 & r_1 & & & \\ & & -r_1 - r_2 & r_1 & r_2 & & & \\ r_3 & & & -r_1 - r_3 & & & r_1 & & \\ & & & & -2r_2 & r_2 & r_2 & & \\ r_3 & & & & -r_1 - r_3 & r_1 & & \\ & & & & & -r_2 - r_3 & r_2 & \\ & & & & & & -r_2 - r_3 & r_2 & \\ & & & & & & & -r_2 - r_3 & r_2 & \\ & & & & & & & -r_2 - r_3 & r_2 & \\ & & & & & & & & -r_2 - r_3 & r_2 & \\ & & & & & & & & -r_2 - r_3 & r_2 & \\ & & & & & & & & -r_2 - r_3 & r_2 & \\ & & & & & & & & -r_2 - r_3 & r_2 & \\ & & & & & & & & & -r_2 - r_3 & r_2 & \\ & & & & & & & & & -2r_3 & \end{pmatrix}$$

The tiny.maple file contains this matrix information in a form suitable for consumption by a solution tool such as Matlab or Maple. When using the Matlab version of the Workbench the transitions file would contain this.

```
Q(1,3) = Q(1,3) + r1;
Q(1,1) = Q(1,1) - r1; % 1 --start,r1-> 3
Q(1,2) = Q(1,2) + r1;
Q(1,1) = Q(1,1) - r1; % 1 --start,r1-> 2
...
```

When using the Maple version of the Workbench the transitions file would contain this.

```
Q[1,3] := Q[1,3] + r1:
Q[1,1] := Q[1,1] - r1: # 1 --start,r1-> 3
Q[1,2] := Q[1,2] + r1:
Q[1,1] := Q[1,1] - r1: # 1 --start,r1-> 2
```

¹The name of this file is instead tiny.m for the Matlab version of the PEPA Workbench.

The essence of the two presentations is the same. A square sparse matrix [the sparse element is zero] is modified by a sequence of assignments which record the rate at which a transition can move the system from one state to another. Explanatory comments, ignored by the solution package, give the identifier of the transition and make clear the direction of the state change. The comments are stating this information: $3 \stackrel{a,r_1}{\longleftrightarrow} 1 \stackrel{a,r_1}{\longrightarrow} 2$. This information is useful to the modeller when examining the state space to detect symmetries within the PEPA model or when 'debugging' a faulty model.

The states of the model have been associated with numbers for use as the indices for the transition matrix. Now we need to understand which state has which number. For this information we look at both the files for the state table tiny.table and for the associated hash table tiny.hash. These contain the information shown in Figure 2. The state table shows the assignment of state numbers to states, identified in terms of their hash table identifier equivalents.

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The st	ate table	The has	sh '	table
1 ⊢	$\rightarrow a \parallel a$	P_1	\mapsto	a
2 ⊢	$\rightarrow d \parallel a$	P_2	\mapsto	d
3 ⊢	$\rightarrow a \parallel d$	P_3	\mapsto	g
4 ⊢	$\rightarrow g \parallel a$	r_1	\mapsto	c
$5 \vdash$	$\rightarrow d \parallel d$	r_2	\mapsto	f
6 ⊢	$\rightarrow a \parallel g$	r_3	\mapsto	i
7 ⊢	$\rightarrow d \parallel g$	run	\mapsto	e
8 ⊢	$\rightarrow g \parallel d$	start	\mapsto	b
9 ⊢	$\rightarrow g \parallel g$	stop	\mapsto	h

Figure 2: State table and hash table for the tiny model.

Why have a state table? The modeller using the Workbench needs to know which states are which because certain states will be distinguished success states or completion states or utilisation states and the experimentation upon the model proceeds by considering the effect on the model of changes of relative rate.

Why have a hash table? All of the limiting problems with the use of the PEPA Workbench when processing large PEPA models are to do with space requirements either in terms of memory usage or disk usage. In generating the state space of a model the Workbench may exhaust the main memory capacity and fail. In comparison the time taken to process a model is of lesser importance because a modeller could simply run the model for longer. Thus when we attempt to improve the PEPA Workbench it is always by seeking to reduce the memory usage [perhaps at the cost of some run-time penalty]. Problems related to state-space size are a concern for all Markov modelling tools.

Notice that in the hash table no letters which were used as identifiers in the PEPA input are re-used as internal identifiers within the Workbench. This will not cause a difficulty if it happens but we take this opportunity to point out that—as with all model development—the user should choose meaningful identifiers for the activities, rates and components of a PEPA model.

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At this point the PEPA Workbench has done its work and it is now time to use the available solution tool to solve the model to find the steady state probability distribution. Performance measures can be calculated from this distribution. For the particular 9×9 matrix generated by the tiny example shown here it is as easy to solve symbolically using Maple (see Appendix B) as it is to solve numerically using Matlab (see Appendix A). For examples of even moderate size seeking symbolic solutions becomes impractical.

PEPA grammar

We now present the input grammar of the ASCII syntax for the PEPA language in EBNF notation.

program	::=	$declaration^+ \ composition$	
declaration	::=	$#id = seq_component$;	
$seq_component$::=	$seq_component / \{ [idseq] \}$	hiding
		seq_component + seq_component	choice
		(id, rate) . seq_component	prefixing
		id	variable
		(seq_component)	grouping
composition	::=	$seq_component < [idseq] >$	
		$seq_component$	
		$seq_component < [idseq] >$	
		composition	
		(composition)	
idseq	::=	id	
		id , idseq	
rate	::=	id	
		int	
		infty	
id	::=	alphanumeric sequence	
int	::=	unsigned numeric sequence	

In addition LATEX-style comments are supported, that is, the PEPA Workbench ignores any characters which follow a percent sign on a line.

The form of the component which follows the declaration sequence should be a parallel composition of (a subset of) the components which have been declared in the declaration sequence.

As noted above, a reserved word is used to represent the infinity symbol which denotes that a component is passive with respect to this action. The reserved word is infty [as in T_EX and IAT_EX] where the symbol \top is used in the mathematical syntax for PEPA.

Example sequential components: If P and Q are the only components declared in the declaration sequence in the PEPA program then the following are legal component expressions which could appear in later declarations.

P /	{a}	Hiding any a action of P
P +	Q	Choice between P or Q
(a,	r1).P	Prefixing P by a with rate r1
(a,	12).P	Prefixing P by a with rate 12
(a,	infty).P	Prefixing ${\tt P}$ by a performed passively

Sequential components non-examples: If P and Q are the only components declared in the declaration sequence in the PEPA program then the following are illegal component expressions for the reason given.

R	Variable not declared
P \ {a}	The hiding symbol is the division sign
P/a	Set brackets are not optional
P <> Q	No use of parallel in sequential components
P.Q	No sequential composition in PEPA
(a, r1) P	Omitted full stop
(a, 12.0).P	Real number constants not allowed
(a, 5 / 2).P	No arithmetic operations; only constants
P[b/a]	No renaming in PEPA

Example parallel compositions: If P and Q are the only components declared in the declaration sequence in the PEPA program then the following are legal parallel compositions which could appear after the declarations.

Ρ	<>	Q							${\tt P}$ and ${\tt Q}$ proceed independently in parallel
Ρ	<a,< td=""><td>, b</td><td>)></td><td>Q</td><td></td><td></td><td></td><td></td><td>${\tt P}$ and ${\tt Q}$ synchronise on ${\tt a}$ and ${\tt b}$</td></a,<>	, b)>	Q					${\tt P}$ and ${\tt Q}$ synchronise on ${\tt a}$ and ${\tt b}$
(F	, <>	> P)	<a,< td=""><td>b></td><td>Q)</td><td><></td><td>Q)</td><td>Copies of $\tt P$ and $\tt Q$ synchronise on $\tt a$ and $\tt b$</td></a,<>	b>	Q)	<>	Q)	Copies of $\tt P$ and $\tt Q$ synchronise on $\tt a$ and $\tt b$

Parallel compositions non-examples: If P and Q are the only components declared in the declaration sequence in the PEPA program then the following are illegal parallel compositions for the reason given.

R	Variable not declared
P Q	Pure parallel composition symbol is $<>$
P < a b > Q	Omitted comma
(a, r1).P	No sequential operations in the composition

Experimenting with a model

Finding the equilibrium probability distribution for a model is an important first step in investigating its behaviour. Once this has been achieved the modeller can conduct a series of experiments to investigate the model thoroughly. The series of experiments which are undertaken will vary from model to model but they commonly involve selecting an interesting subset of the state space and finding the probability of being in those states. An activity which must be performed in doing this part of the experimentation is finding the numbers which are associated with the states which are of interest. We provide another tool to assist with this: the tool is called the PEPA State Finder.

The modeller first describes some states of interest through the use of a simple pattern language with stars for wild cards and vertical bars for separators between model components. Returning to our tiny example, a modeller interested in testing how often either component in the tiny example was component P1 would prepare a file called tiny.psf with the following contents.

```
% We are checking for P1 in either case
test: P1 |*
test: * |P1
```

The identifier test gives the name of the function and after the colon comes the pattern of interest. Assuming that you access the PEPA State Finder by typing **psf** then the function is processed as shown in Figure 3 where the user only typed the **psf** command, the model name and the function name and the rest was produced by the tool.

```
[unix]: psf
PEPA State Finder [Version 0.06, solver, 2-12-1997]
PEPA model name: tiny
PSF file name: tiny
Function ''test'' written to file ''test.fun''.
Exiting PEPA State Finder.
```

Figure 3: A sample PEPA State Finder session

The function which was generated by this is shown in Appendix C. For this tiny example the function could easily have been created by hand since the model has a very regular structure and only nine states but with larger models the usefulness of the PEPA State Finder becomes clearer. Notice that here the tool correctly avoids counting the state $P_1 \parallel P_1$ twice. This would be a mistake which could easily be made otherwise.

The PEPA State Finder must be run after running the PEPA Workbench on a model because it searches through the table file which the Workbench produces.

Limitations

The PEPA Workbench has a number of limitations.

- The Workbench implements Cyclic PEPA only (see [Hillston, 1996] for the definition of this term). These are the only PEPA models for which steady-state probability distributions can be calculated. This is a consequence of the language definition and will not be changed.
- The Workbench does not implement the apparent rates of PEPA. All synchronisations must be between one active component and one passive component. This is an error which should be corrected in later versions of the PEPA Workbench.

Appendix A: A sample MATLAB session

In this appendix the nine state model is solved for the particular case when $r_1 = r_2 = r_3 = 2.0$. Since all of the transitions proceed at the same rate, and the model is symmetric, all of the states are equally likely, with calculated probability $\frac{1}{9} \approx 0.1111$.

```
[unix]: matlab
```

< M A T L A B (R) >
(c) Copyright 1984-94 The MathWorks, Inc.
 All Rights Reserved
 Version 4.2c
 Dec 31 1994

Commands to get started: intro, demo, help help Commands for more information: help, whatsnew, info, subscribe

```
>> Q=sparse(9,9);
>> r1=2.0;
>> r2=2.0;
>> r3=2.0;
>> tiny;
>> for i=1:9
        Q(i,9) = 1.0;
end
>> b=zeros(9,1);
>> b(9) = 1.0;
>> QT = Q';
>> P = QT \ b
```

```
P =
```

0.1111 0.1111 0.1111 0.1111 0.1111 0.1111 0.1111 0.1111 0.1111

>> quit

1020 flops.

Appendix B: A sample Maple session

In this appendix the nine state model is solved symbolically for all values of r_1 , r_2 and r_3 .

```
[unix]: maple
  |\^/|
         Maple V Release 3 (The University of Edinburgh)
_//| //_ Copyright (c) 1981-1994 by Waterloo Maple Software and the
\ MAPLE / University of Waterloo. All rights reserved. Maple and Maple V
<____> are registered trademarks of Waterloo Maple Software.
          Type ? for help.
    > with(linalq):
Warning: new definition for
                           norm
Warning: new definition for
                           trace
> Q := array(sparse, 1..9, 1..9):
> read 'tiny.maple':
> b := array(sparse, 1..9):
> for i to 9 do
     Q[i,9] := 1.0
 od:
> b[9] := 1:
> QT := transpose(Q):
> P := linsolve(QT,b);
           2 2 2
                            2
                                        2
                                              2 2
          r3 r2 r3 r2 r1 r3 r2 r1 r3 r2 r1 r3 r1
    P := [ -----, -----, -----, -----, -----,
                           %1
            %1 %1
                                     %1
                                                %1
                             2 2 2
                      2
           2
       r3 r2 r1 r3 r2 r1 r3 r2 r1 r1 r2
       -----, ------, ------, ------ ]
          %1 %1
                           %1
                                   %1
                 2 2
           2
                          2 2
                                         2
%1 := 2. r3 r2 r1 + r3 r2 + r1 r2 + 2. r2 r1 r3 +
             2 2 2
     2. r2 r1 r3 + r3 r1
> quit
bytes used=983796, alloc=786288, time=1.42
```

Appendix C: A function generated by the PEPA state finder

This is the function which was generated when the tiny.psf file was processed by the PEPA State Finder. The body of the test function adds together the probabilities for the relevant states. After a comment symbol on each line comes the description of the state. As we expected, all of the states contain a P_1 component.