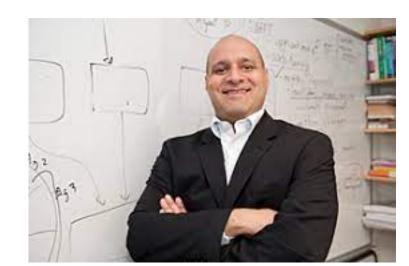
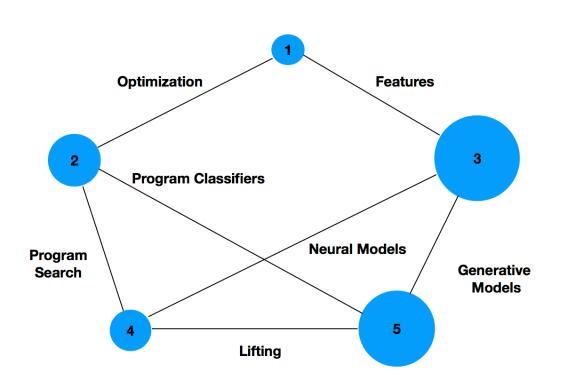
# Rethinking Compilation: L5







| Program |  | →x86 — | → Hardware |
|---------|--|--------|------------|
|         |  |        |            |

Program — OpenCL — Hardware

Program → clBLAS — Hardware Hardware

# Overview<sup>1</sup>

- L1: Motivation and survey of auto-tuning/machine learning for compilers
- L2: Program rewriting schemes e-graphs and equality saturation
- L3: Program embeddings and Graph Neural Networks
- L4: Program synthesis and neural synthesis
- L5: Neural Machine Translation, Transformers and Large language models

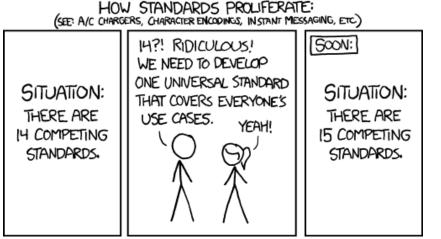
# Lecture Structure

- 1. The Tower of Babel of Programming Languages
- 2. Machine Translation of Programming Languages
- 3. Transformers
- 4. Unsupervised Translation
- 5. Translation validation
- 6. What's next

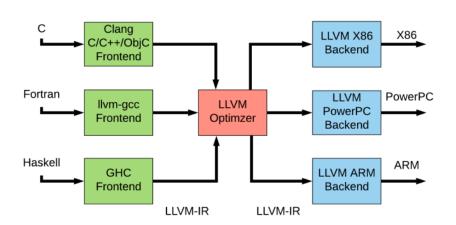
#### The Tower of Babel of Programming Languages

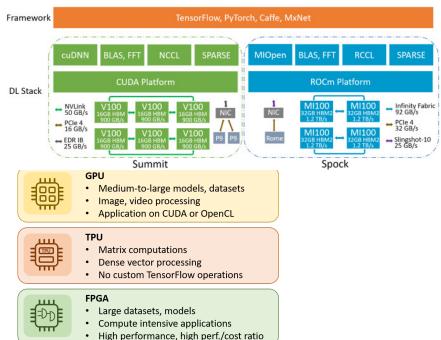
# Ultimate language/compiler/backend library to rule them all!

https://xkcd.com/927/



#### The Tower of Babel of Programming Languages





#### The Tower of Babel of Programming languages: Automation?

#### A: Translate/rewrite

- generate x86 code from LLVM IR, or remove dead code
- **B:** Decide when to rewrite given (A)
- Vectorize a loop.

Large potential for automation. ML in compilers used in B, but A

Replacing heuristics is an easy, controllable win for ML

Generating code is way more challenging and dangerous!

#### The Tower of Babel of Programming Languages

Ways forward to translation automation

- A. Compilers, developer tools.
- B. Program synthesis.
- C. Proof automation, proof repair.
- **D.** Machine translation?

- + Reliable,
- Less flexible, require manual effort
- + **Flexible**, potentially fully automated
- Less reliable!

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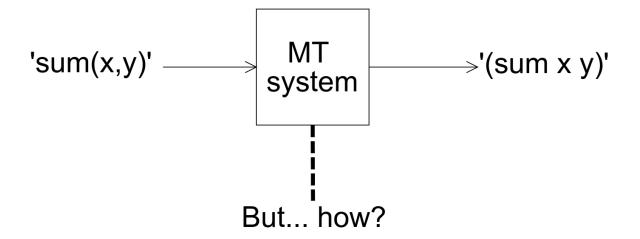
Translate language S into language T.

Typically, instances in S and T are expressed as discrete **sequences of tokens**.

$$s = [s_1, \dots s_n] \in S, t = [t_1, \dots t_m] \in T$$

Originally developed for natural languages

Potentially applicable to any **discrete domain!** 



# Machine translation for natural languages:

- A. Rule-based MT: Grammars, dictionaries... PL-like approaches
- B. Statistical MT: Faster and less data-hungry than NMT

#### C. Neural MT:

- a. Neural networks feature extractors
- b. Ranking translations
- **c. End-to-end** (i.e. let the model do everything)

Focus on end-to-end NMT.

Tokenization: Assuming a finite vocabulary {'x': 0, '(': 1, ')': 2, ',': 3, 'y': 4, 'sum': 5}

$$sum(x,y)'$$
 -> [sum', '(', 'x', ',', 'y', ')'] -> [5, 1, 0, 3, 4, 2]

Avoid **out-of-vocabulary** words with **finite** vocabulary?

"my\_variable\_2" identifier

Solution: **subword encoding** 

#### Subword encoding:

## Training:

- 1. Initialize vocabulary with all characters
- 2. Merge most frequent pairs
- 3. Until desired vocabulary size reached.

#### Inference:

- 1. If token in vocabulary done
- 2. Otherwise, express with subtokens

List of token indices: [5, 1, 0, 3, 4, 2]

Neural networks need continuous vectors.

Solution: introduce embedding lookup table.

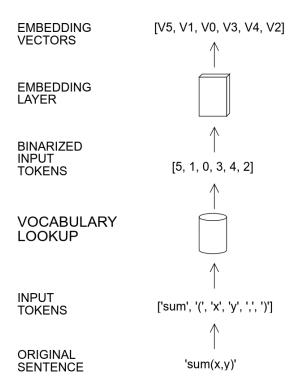
0: [0.754, -0.25, ...]

1: [-0.11, ...]

2: ...

**Initialized** to **random** values **Optimized end to end** with the rest of the network.

Vectors will get closer to the vector that minimizes the error of predictions



vectors of floats (NN-friendly).

But variable-length list of them (NN-unfriendly).

multi-layer perceptron (MLP) requires a fixed size vector input.

#### Several possibilities:

- Apply an MLP to each vector, and aggregate all outputs somehow (mean? sum?)
- Recurrent neural network LSTM, a well-known RNN architecture
- Other

## **Seq2seq**, an **encoder-decoder** architecture.

- Encoder architecture: RNN (Recurrent Neural Network).
- Decoder architecture: another RNN.

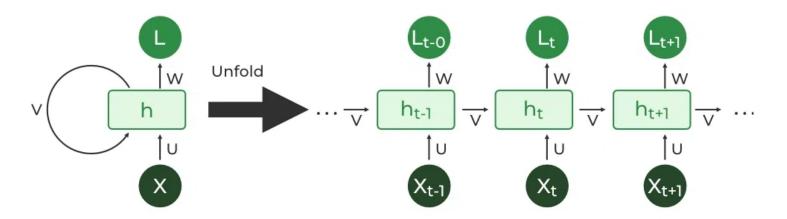
#### **Encoder:**

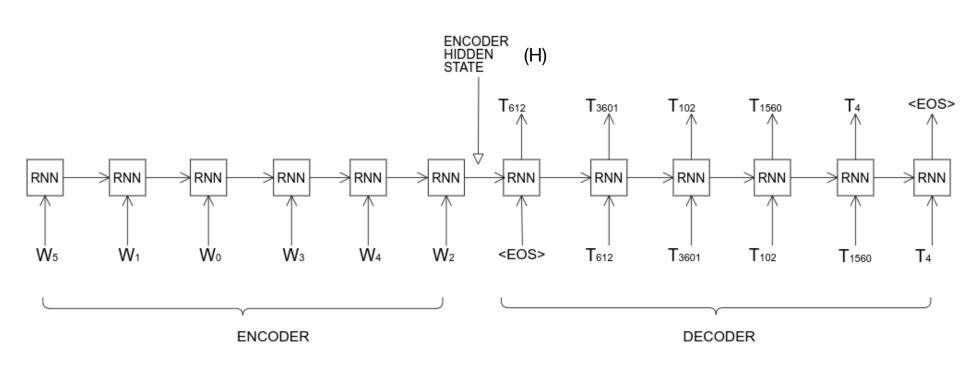
- variable-length sequence of embedding vectors
- into a single, fixed-size vector.

**Decoder**: fixed-size vector (from the encoder) into a variable-length sequence of target token indices.

#### RNN for each token

feed the embedding of the token itself and the RNN output from the previous token (state/"memory")

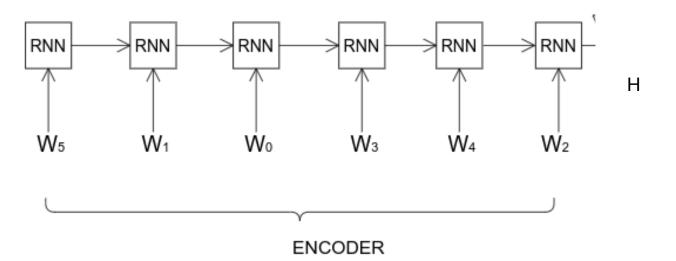




Encoder(source tokens):

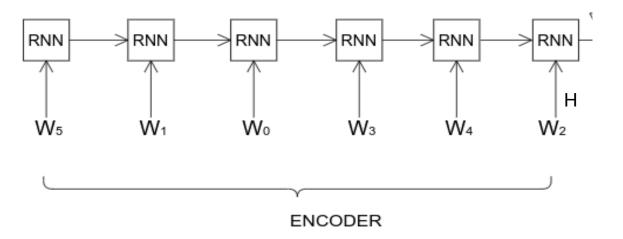
- returns single, fixed-size vector **H** ("hidden state") representing the source sequence.

It's the output of the RNN after the last token!



We have transformed the **original source string into a single vector of floats** (H)!

What do we do now with this H?



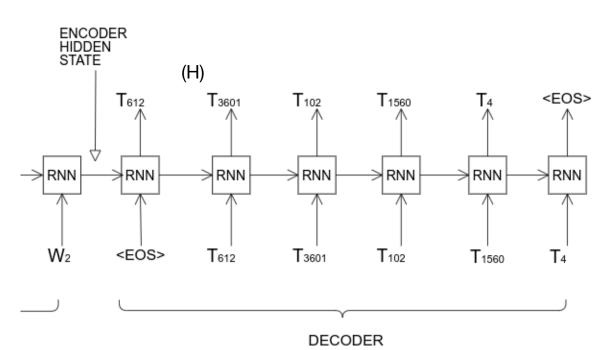
Decoder(H, current\_prediction=[]):

Returns next target token.

Decoder, also an RNN.

# But **not the same RNN!**

It has independent parameters.

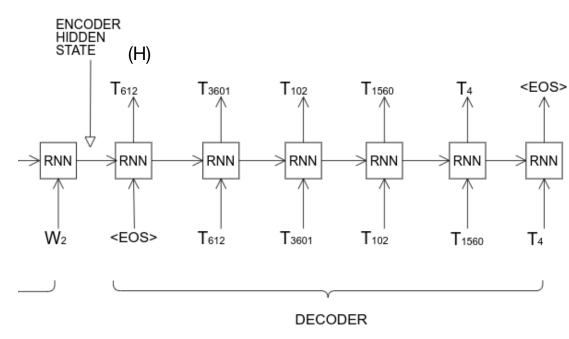


# Second "time step" (**T3601**) decoder reads:

- Encoder final output H.
- Previous Decoder output

#### Output:

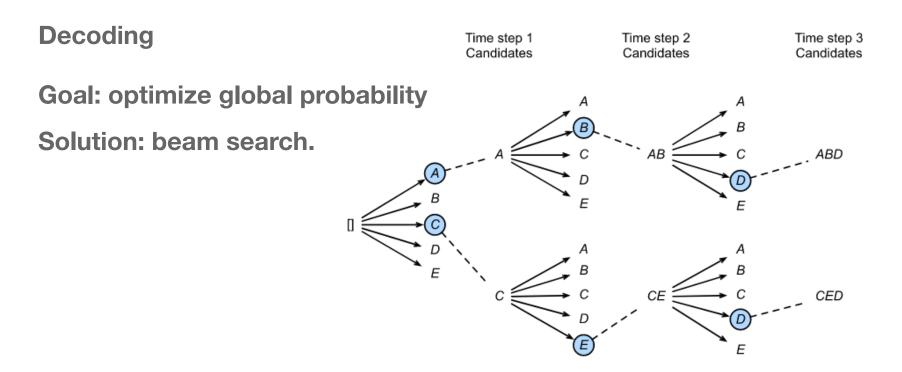
- Useful representation for the next time step.
- Probability over the target vocabulary,
- T3601 has a very high probability



**Decoding**: inference predictions out of the decoder

## Simplest approach: greedy decoding.

- Argmax for each individual token prediction.
- Problem: maximizes local probability, not the global one!

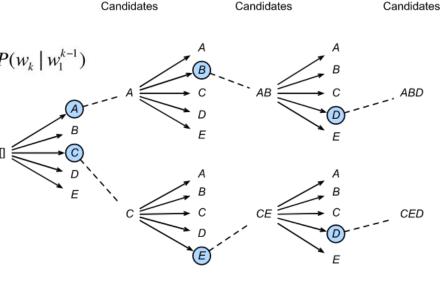


Probability of the sequence:

 $P(w_1^n) = P(w_1)P(w_2 \mid w_1)P(w_3 \mid w_1^2)...P(w_n \mid w_1^{n-1}) = \prod_{k=1}^n P(w_k \mid w_1^{k-1})$ 

Beam search with beam size=k

We keep the top k most probable sentences along all paths.



Time step 2

Time step 1

Time step 3

#### The information bottleneck:

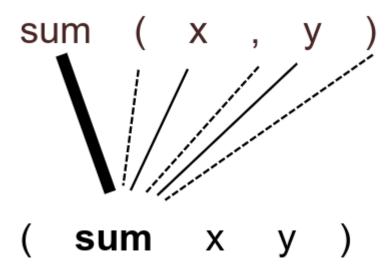
Encode the source with a single, fixed-size vector without losing information?

Solution: attention, i.e. weighted average over source hidden states.

Weights are dynamic (depend on the current decoder state).

Learned end-to-end with the rest of the neural network

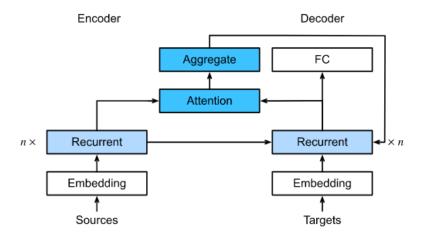
#### Attention



Bahdanau attention: Seq2seq ++

Encoder identical to Seq2seq.

Weighted average of all source hidden states. The weights depend on each decoder state.



# Can we do better than that? Transformers

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# **Transformers**

Transformers are eating the world [4]

#### **Attention Is All You Need**

Ashish Vaswani\* Google Brain avaswani@google.com

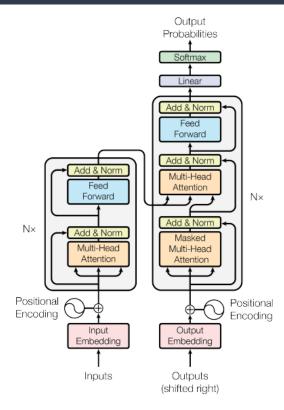
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Noam Shazeer\* Google Brain noam@google.com Niki Parmar\* Google Research nikip@google.com Jakob Uszkoreit\* Google Research usz@google.com

Aidan N. Gomez\* † University of Toronto aidan@cs.toronto.edu Łukasz Kaiser\* Google Brain lukaszkaiser@google.com

Illia Polosukhin\* † illia.polosukhin@gmail.com

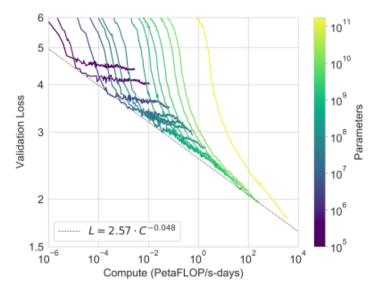


#### **Transformers**

Especially in the form of large language models!

#### **Language Models are Few-Shot Learners**

| Tom B. Brov               | wn* Benjan        | nin Mann*   | Nick R    | kyder* Mel         | lanie Subbiah* |
|---------------------------|-------------------|-------------|-----------|--------------------|----------------|
| Jared Kaplan <sup>†</sup> | Prafulla Dhariwal | Arvind Ne   | elakantan | Pranav Shyam       | Girish Sastry  |
| Amanda Askell             | Sandhini Agarwal  | Ariel Herbe | rt-Voss   | Gretchen Krueger   | Tom Henighan   |
| Rewon Child               | Aditya Ramesh     | Daniel M.   | Ziegler   | Jeffrey Wu         | Clemens Winter |
| Christopher He            | sse Mark Ch       | en Eric S   | igler     | Mateusz Litwin     | Scott Gray     |
| Benjamin Chess            |                   | Jack Clar   | k         | Christopher Berner |                |
| Sam McCandlish Alec Ra    |                   | Radford     | Ilya Su   | itskever I         | Dario Amodei   |



35

#### **Transformers**

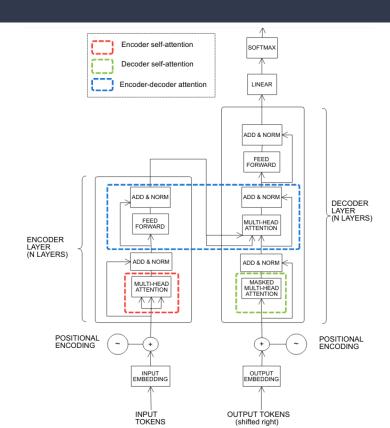
**Transformers** "all-attention" **encoder-decoder** NN (**no recurrence**).

High-level - similar Seq2seq.

Low-level - different

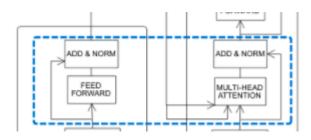
• parallel, pair-wise attention instead of sequential, recurrent network

## **Transformers**



Seq2seq++ **encoder-decoder attention**, attend to the source sequence wrt the target sequence.

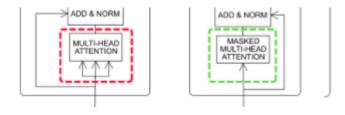
The Transformer **also** has encoder-decoder attention, no fundamental difference here!



Implementation is slightly more complicated than Bahdanau attention.

#### Replaces RNNs with self-attention

- in the encoder and the decoder
- attend to the source sequence wrt the source sequence itself, in the encoder,
- attend to the target sequence wrt the target sequence itself, in the decoder.



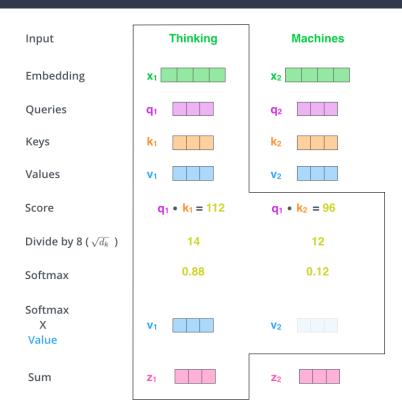
Attention implementation: scaled dot-product attention

Queries: "Is X information present?"

**Keys**: "Is information present in Y word?"

Use query value dot product as the importance weight in the weighted average.

**Value**: "Now that we know the weight importance, what's the information that we wanted to transfer?"



All attention implemented identically.

Difference is where queries, keys, and values come from:

- Encoder: the source sequence.
- Decoder: the target sequence.
- Encoder-decoder: Queries come from the target sequence. Keys and Values come from the source sequence.

Encoder attention can be applied in parallel

Decoder cannot run in parallel in inference

Transformers scale better than previously thought.

Quadratic complexity with respect to the sequence length n.

#### But

- As **embedding dimension d increase.** n becomes less important.
- Some tasks (??) can be solved in **relatively small sequence length n**.

| Layer Type     | Complexity per Layer       | Sequential Operations | Maximum Path Length |
|----------------|----------------------------|-----------------------|---------------------|
| Self-Attention | $O(n^2 \times d)$          | O(1)                  | O(1)                |
| Recurrent      | $O(n \times d^2)$          | O(n)                  | O(n)                |
| Convolutional  | $O(k \times n \times d^2)$ | O(1)                  | $O(\log_k(n))$      |

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NMT large parallel datasets S, T.

In NLP pairs (S,T) with the same meaning.

- Subtitled movies or European Parliament transcriptions
- Using **sentence alignment** algorithms.

#### PL more difficult:

- need function-to-function parallel datasets
- no natural occurrences of them (Exception: Code-assembly pairs).

Unsupervised translation more important in programming languages!

### Unsupervised learning

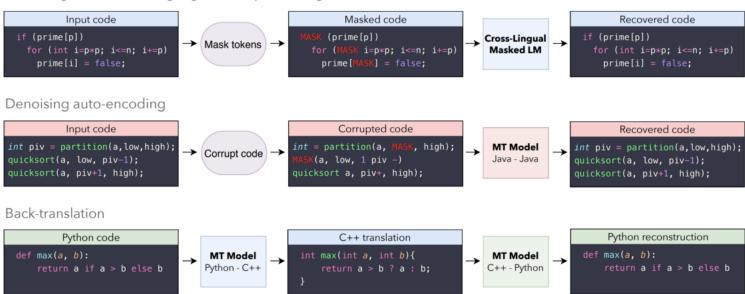
- Unsupervised algorithms?
- Apply supervised algorithms to unlabelled data.

#### Basic idea:

- automatically **create a low-quality, synthetic parallel dataset**
- bootstrap the learning.
- Iteratively improve.

# Unsupervised Translation of Programming Languages (Transcoder)[7]

Cross-lingual Masked Language Model pretraining



### Create a synthetic dataset.

• Step 1: Pretrain an encoder

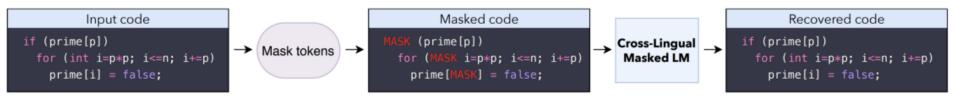
#### Predict masked tokens.

E.g. "def <MASK>(x,y): return x + y" -> predict "sum".

Randomly replace tokens from the functions with a special token <MASK>.

- Train the encoder to predict the masked tokens
- (Similar to word2vec in L3)

#### Cross-lingual Masked Language Model pretraining



#### Randomly mask functions (**Unsupervised**).

Pretraining task forces the encoder to learn useful embeddings

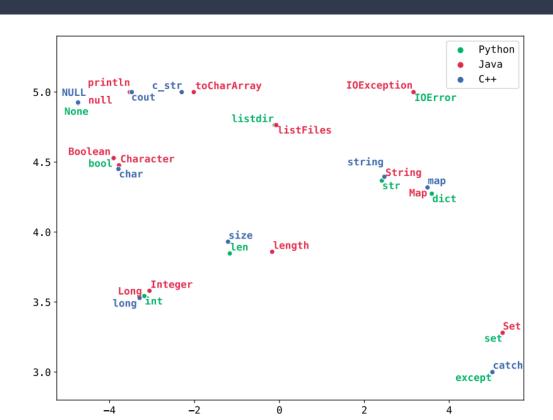
### Learn aligned embeddings,

 same encoder to simultaneously demask functions in language A and B shared vocabulary and embedding table.

### Same model for both,

- forces alignment
- tokens in A occur in semantically similar contexts in B

Tokens used in similar contexts cluster together even if they belong to different languages.



Step 2: Use pre-trained, multilingual encoder to **initialize an encoder-decoder Transformer**.

Encoder identical to the pretrained one.

Decoder identical to the encoder

except randomly initialize the cross-attention modules

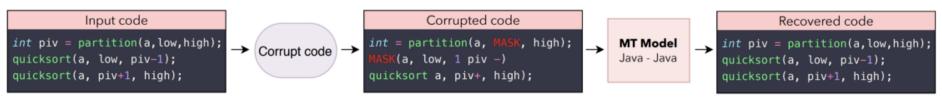
Train this encoder-decoder model to **demask** the original sentence

But using **special tokens** to denote the desired target (e.g., <TO\_PYTHON>).

### Example:

"def **<MASK>** (x, y): x+y" -> encoder + **<TO\_PYTHON>** -> "def **sum**(x, y): x+y"

#### Denoising auto-encoding



Step 3: **Zero-shot translation** (unseen directions).

Trained Python to Python and Java to Java.

After feeding a Java function can we inject <TO\_PYTHON> instead of <TO\_JAVA> ?

Trained the same model on Java and on Python.

Can we translate without having done it in training?

#### Yes the embeddings are aligned

the model is going to extrapolate to perform a poor-quality translation!

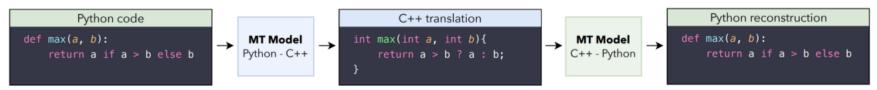
Poor-quality Java-Python translations are not useful on their own

But can build a parallel synthetic dataset

For every Python sentence:

- 1. Apply (bad) translator to get C++ -ish
- 2. Use the reverse direction to build a synthetic dataset **back translation**

#### Back-translation



Use synthetic dataset to learn translator using supervised learning.

### iteratively improve

Use the new translator (better than the zero-shot one)

Better synthetic dataset and train an even better translator.

## Unsupervised translation will be key for PL/compilers

Cross-lingual Masked Language Model pretraining Input code Masked code Recovered code if (prime[p]) (prime[p]) Cross-Lingual if (prime[p]) Mask tokens → Masked LM for (int i=p\*p; i<=n; i+=p) for (MASK i=p\*p; i<=n; i+=p)</pre> for (int i=p\*p; i<=n; i+=p)</pre> prime[i] = false; prime[MASK] = false; prime[i] = false; Denoising auto-encoding Input code Corrupted code Recovered code int piv = partition(a,low,high); int piv = partition(a,low,high); int = partition(a, MASK, high); MT Model Corrupt code → Java - Java quicksort(a, low, piv-1); (a, low, 1 piv -) quicksort(a, low, piv-1); quicksort(a, piv+1, high); quicksort a, piv+, high); quicksort(a, piv+1, high); Back-translation Python code C++ translation Python reconstruction def max(a, b): def max(a, b): int max(int a, int b){ MT Model MT Model C++ - Python return a if a > b else b return a if a > b else b Python - C++ return a > b ? a : b;

```
S = deque()
                                                             deque <int> S;
G = deque()
                                                             deque <int> G;
for i in range(k):
                                                            for(int i = 0; i < k; i ++){
  while (len(S) > 0 \text{ and } arr[S[-1]] >= arr[i]):
                                                               while((int) S.size() > 0 && arr[S.back()] >= arr[i])
    S.pop()
                                                                 S.pop_back();
  while (len(G) > 0 \text{ and } arr[G[-1]] \le arr[i]):
                                                               while((int) G.size() > 0 && arr[G.back()] <= arr[i])</pre>
    G.pop()
                                                                 G.pop_back();
  G.append(i)
                                                              G.push_back(i);
                                                               S.push_back(i);
  S.append(i)
for i in range(k, n):
  Sum += arr[S[0]] + arr[G[0]]
                                                            for(int i = k; i < n; i ++){
  while (len(S) > 0 \text{ and } S[0] \le i - k):
                                                               Sum += arr[S.front()] + arr[G.front()];
                                                               while((int) S.size() > 0 && S.front() <= i - k)</pre>
    S.popleft()
  while (len(G) > 0 \text{ and } G[0] \le i - k):
                                                                 S.pop_front();
                                                               while((int) G.size() > 0 && G.front() <= i - k)</pre>
    G.popleft()
  while (len(S) > 0 \text{ and } arr[S[-1]] >= arr[i]):
                                                                 G.pop_front();
    S.pop()
                                                               while((int) S.size() > 0 && arr[S.back()] >= arr[i])
  while (len(G) > 0 \text{ and } arr[G[-1]] \le arr[i]):
                                                                 S.pop_back();
                                                               while((int) G.size() > 0 && arr[G.back()] <= arr[i])</pre>
    G.pop()
  G.append(i)
                                                                 G.pop_back();
  S.append(i)
                                                              G.push_back(i);
Sum += arr[S[0]] + arr[G[0]]
                                                               S.push_back(i);
return Sum
                                                             Sum += arr[S.front()] + arr[G.front()];
                                                            return Sum;
24 to 60% accuracy
```

int Sum = 0;

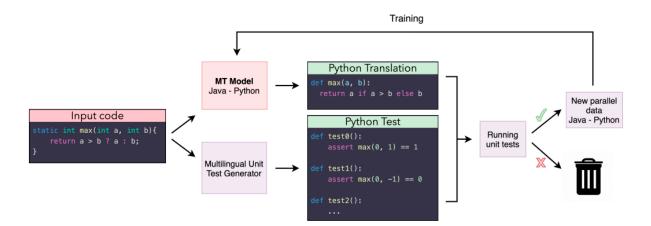
int SumOfKsubArray(int arr[], int n, int k){

def SumOfKsubArray(arr, n, k):

Sum = 0

Bonus: select the synthetic functions

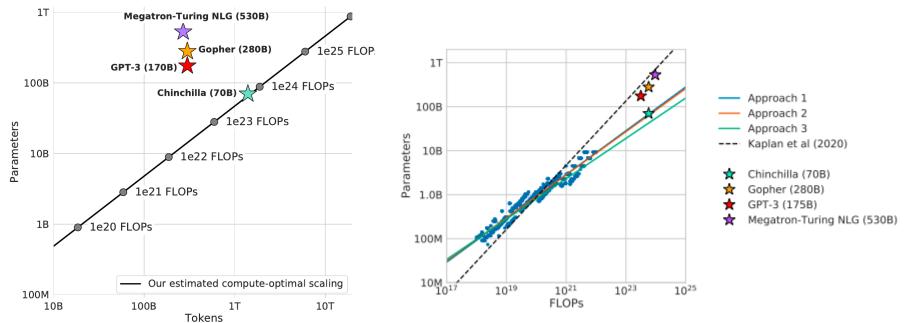
- discard the bad ones by running the original unit tests on them.
- discard the syntactically incorrect synthetic functions.



#### Language models

Neural **Scaling Laws**: neural language models scale smoothly with optimally allocated compute.

No diminishing returns yet?!



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Evaluating text generation hard

Natural languages, require human evaluation

- slow and expensive.

Solution: **similarity** metrics **correlate** with human evaluation.

BLEU (n-gram match) in academia and edit similarity in the industry.

#### **BLEU** in Action

(Source Original) 枪手被警方击毙.

(Reference Translation) The gunman was shot to death by the police



What about code translation?

Disadvantage: hard/binary correctness, a single character may corrupt an otherwise perfect solution.

Advantage: well-defined syntax and semantics

- Exact match: wrt ground truth translation.
- CodeBLEU: BLEU with AST features.
- IO/unit tests: correct iff passes unit tests ("observational equivalence")
- Formally verified: equivalent to source sequence.

## Compiler validation vs translator validation:

- Compilers/synthesizers are correct by design/construction
- Translator validation >>> compiler validation
- Growing interest in neural network interpretability/verifiability

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## What's next

Methods are generic: apply to any sequence

Add inductive biases from PL?

- graph neural networks. Good at code classification, but underperform Transformers at generative tasks!
- "bitter lesson" for ML for code: "just" scaling up outperforms everything?

Can we learn from program synthesis and compilers?

## What's next

Code inductive biases are likely to inform future designs

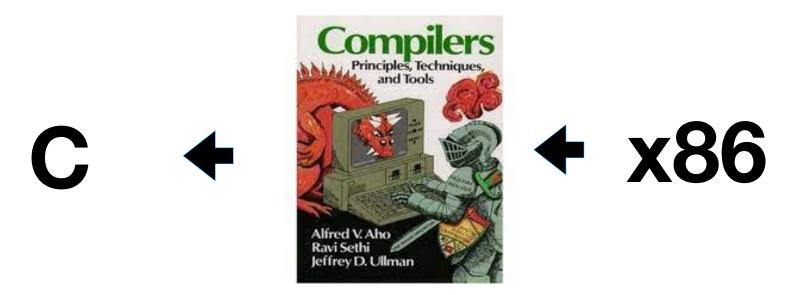
Best changes will not necessarily be in the architecture level:

- Data: can we preprocess data differently/better?
- Training objectives: can we find better loss functions for the code domain?
- Hybrid methods: can we effectively combine NMT with program synthesis?
- Input/output pairs: can we use input/output pairs to guide the translation?

## David vs Goliath: Decompilation and ChatGPT

(SLaDe: A Portable Small Language Model Decompiler for Optimized Assembler)

# Decompilation: x86->C

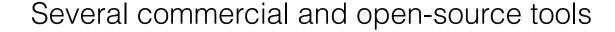


Difficult to run compilers backwards

# Decompilation: x86->C

## Explored for 50 years

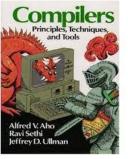
- Used as precursor for binary translation
- Surprising lack of evaluation on correctness
- Often produce mangled code



Many person-years of effort

- Retdec over 25 person-years







## State-of-the-art GHIDRA



Decompiler from NSA Many person-years effort



#### Useful for:

- Security purposes!
- Porting legacy code.
- Lifting?

Excellent rule-based decompiler

... but produces hard-to-read code!



What does this code do?

### GCC 03:

```
.globl add
 .type add, @function
add:
                         .L3:
.LFB0:
 .cfi_startproc
                         movslq %edi, %rax
endbr64
                         leal 1(%rdi), %r8d
movq %rdi, %rcx
                          salq $2, %rax
testl %edx, %edx
                          addl %esi, (%rcx, %rax)
jle .L1
leal -1(%rdx), %eax
                          cmpl %r8d, %edx
cmpl $2, %eax
                         jle .L1
jbe .L6
                         addl $2, %edi
movq %rdi, %rax
                          addl %esi, 4(%rcx,%rax)
movl %edx, %edi
movd %esi, %xmm2
                         cmpl %edi, %edx
shrl $2, %edi
                         jle .L1
pshufd $0, %xmm2,
                          addl %esi, 8(%rcx,%rax)
       %xmm1
                         .L1:
subl $1, %edi
salq $4, %rdi
                         ret
leaq 16(%rcx, %rdi), %rdi
                         .p2align 4,,10
 .p2align 4,,10
                          .p2align 3
 .p2align 3
                         .L9:
.L4:
                         ret
movdqu (%rax), %xmm0
addq $16, %rax
                         .L6:
paddd %xmm1, %xmm0
                         xorl %edi, %edi
movups %xmm0, -16(%rax)
                         jmp .L3
cmpq %rdi, %rax
                          .cfi_endproc
jne .L4
movl %edx, %edi
andl $-4, %edi
testb $3, %dl
je .L9
```



Generates structured code

- but difficult to read

Non-intuitive variable names

Use of shifts and masks

Follows O3 control-flow mangling

```
void add(int *param 1, int param 2,
        uint param 3) {
  uint uVar1:
 int *piVar2;
 int *piVar3;
 if (0 < (int)param 3) {
   if (param 3 - 1 < 3) {
      uVar1 = 0;
      do {
        param 1[(int)uVar1] =
          param_1[(int)uVar1] + param_2;
        uVar1 = uVar1 + 1;
      } while ((int)uVar1 < (int)param_3);</pre>
      return:
    } else {
      piVar2 = param 1;
        piVar3 = piVar2 + 4;
        *piVar2 = *piVar2 + param_2;
        piVar2[1] = piVar2[1] + param_2;
        piVar2[2] = piVar2[2] + param_2;
        piVar2[3] = piVar2[3] + param_2;
        piVar2 = piVar3;
     } while (piVar3 = param_1 +
        ((ulong)param 3 >> 2)* 4,
        piVar2 != piVar3);
     uVar1 = param 3 &
        Oxfffffffc;
      if ((param 3 & 3) == 0) {
        return;
      param 1[(int)uVar1] =
        param_1[(int)uVar1] + param_2;
      if ((int)(uVar1 + 1) <
          (int)param_3) {
        param 1[(int)(uVar1 + 1)] =
          param_1[(int)(uVar1 + 1)] +
          param 2:
        if ((int)(uVar1 + 2) <
            (int)param 3) {
          param 1[(int)(uVar1 + 2)] =
            param_1[(int)(uVar1 + 2)] +
            param 2;
          return;
```



```
// Add the aligned elements
                                                    int count = (n - pre_count) / 4;
                                                    m128i* sse src =
void add(int* dest, int* src, int n) {
                                                         (__m128i*) (src + pre_count);
  if (n \le 0) {
                                                   m128i* sse dest =
    return;
                                                         ( m128i*) (aligned dest);
                                                   m128i \times mm1 =
  // If n is smaller than 2, use a loop
                                                         mm loadu si128(sse src);
  // to add the values
                                                   for (int i = 0; i < count; i++) {
  if (n < 2) {
                                                     m128i \times mm0 =
    for (int i = 0; i < n; i++) {
                                                       _mm_load_si128(sse_dest);
      dest[i] += src[i];
                                                     xmm0 = _mm_add_epi32(xmm0, xmm1);
                                                       _mm_store_si128(sse_dest, xmm0);
    return;
                                                     sse src++;
                                                     sse_dest++;
  // Align dest pointer to
  // 16-byte boundary
  int* aligned_dest = (int*)
                                                   // Add the remaining elements
     (((uintptr_t) dest + 15) & ~15);
                                                   int remaining = n -
                                                      (pre_count + count * 4);
                                                   for (int i = 0; i < remaining; i++) {
  // Calculate the number of elements
                                                      dest[pre_count + count * 4 + i] +=
  // before the aligned dest pointer
  int pre_count = (aligned_dest - dest);
                                                         src[pre_count + count * 4 + i];
  // Add the unaligned elements
  for (int i = 0; i < pre_count; i++) {
    dest[i] += src[i];
```



### Compared to Ghidra

- + More readable code than Ghidra
- + More meaningful variable names
- + Even adds comments!
- + Compiles and executes

#### However

- complex code
- introduces x86 intrinsics
- gives incorrect results

```
void add(int* dest, int* src, int n) {
 if (n \le 0) {
   return:
 // If n is smaller than 2, use a loop
 // to add the values
   for (int i = 0; i < n; i++) {
     dest[i] += src[i];
 // Align dest pointer to
 // 16-byte boundary
 int* aligned dest = (int*)
     (((uintptr_t) dest + 15) & ~15);
 // Calculate the number of elements
 // before the aligned dest pointer
 int pre_count = (aligned_dest - dest)
 // Add the unaligned elements
 for (int i = 0; i < pre_count; i++) {
   dest[i] += src[i];
 // Add the aligned elements
 int count = (n - pre_count) / 4;
  m128i* sse src =
        (__m128i*) (src + pre count);
  m128i* sse dest =
        ( m128i*) (aligned dest);
  m128i xmm1 =
        mm_loadu_si128(sse_src);
 for (int i = 0; i < count; i++) {
    m128i xmm0 =
      mm load si128(sse dest);
   xmm0 = mm_add_epi32(xmm0, xmm1);
      mm_store_si128(sse_dest, xmm0);
   sse src++;
   sse dest++;
 // Add the remaining elements
 int remaining = n -
    (pre_count + count * 4);
 for (int i = 0; i < remaining; i++) {
   dest[pre count + count * 4 + i] +=
      src[pre count + count * 4 + i];
```

### SLaDe: both correct and readable!

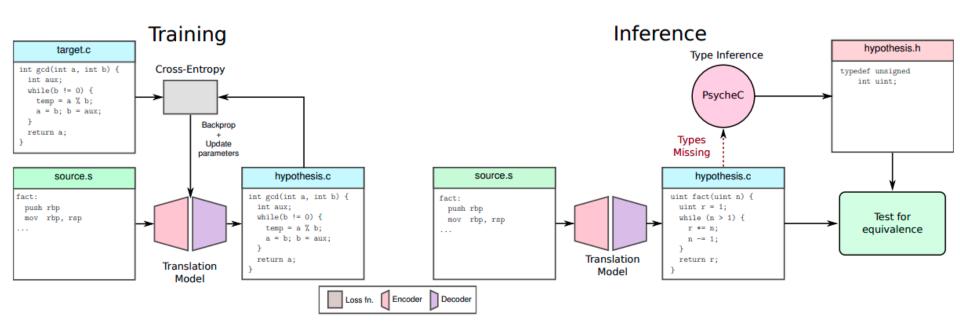
### **Original Source**

```
void add(int *list, int val, int n) {
  int i;
  for (i = 0; i < n; ++i) {
    list[i] += val;
  }
}</pre>
```

```
void add(int a[], int x, int y) {
   int i;
   for (i = 0; i < y; i++) {
      a[i] += x;
   }
}</pre>
```

SLaDe

## SLaDe architecture: small Transformer + Type Inference



## Datasets: AnghaBench (large-scale training and evaluation)

# ANGHABENCH: a Suite with One Million Compilable C Benchmarks for Code-Size Reduction

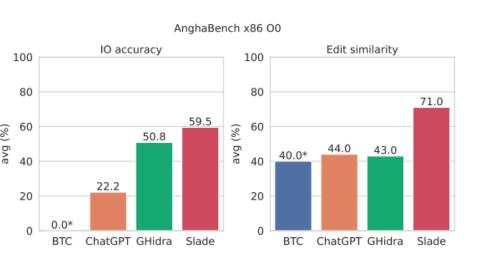
Anderson Faustino da Silva Department of Informatics UEM, Brazil anderson@din.uem.br

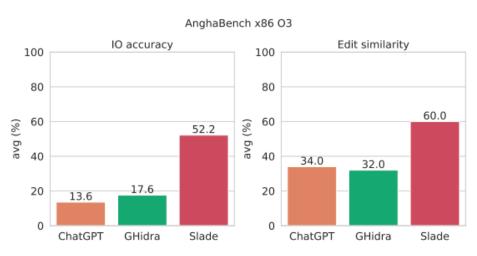
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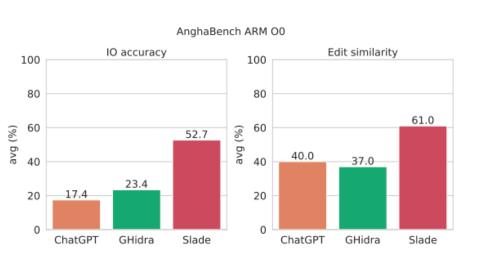
Fernando Magno Quintão Pereira Department of Computer Science UFMG, Brazil fernando@dcc.ufmg.br

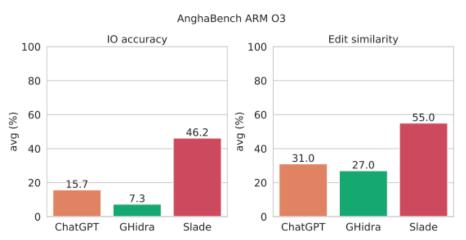
# x86: 3x improvement on O3





# ARM: 3x to 6x improvement on O3





## Analysis

#### Ghidra

- code complexity causes problems
- fundamentally cannot predict types of external functions

#### ChatGPT

- performs well on x86 O0 but O3 causes significant difficulty
- worryingly produces compilable code that is wrong

#### SLaDe

- rarely produces compilable code that is incorrect
- can be improved with program analysis interacting with decoder

### Conclusion

To adapt to a world of language/hardware innovation

We need to rethink compilation

Lots of unusual technology - great time for research!

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# **Rethinking Compilation**

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